# A Simple But Effective Canonical Dual Theory Unified Algorithm for Global Optimization 

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#### Abstract

Numerical global optimization methods are often very time consuming and could not be applied for high-dimensional nonconvex/nonsmooth optimization problems. Due to the nonconvexity/nonsmoothness, directly solving the primal problems sometimes is very difficult. This paper presents a very simple but very effective canonical duality theory (CDT) unified global optimization algorithm. This algorithm has convergence is proved in this paper. More important, for this CDT-unified algorithm, numerous numerical computational results show that it is very powerful not only for solving low-dimensional but also for solving highdimensional nonconvex/nonsmooth optimization problems, and the global optimal solutions can be easily and elegantly got with zero dual gap.


## 1 Introduction

In recent years large-scale global optimization (GO) problems have drawn considerable attention. These problems have many applications, in particular in data mining, computational biology/chemistry. Numerical methods for GO are often very time consum-ing and could not be applied for high-dimensional nonconvex/nonsmooth optimization problems. Due to the nonconvexity/nonsmoothness, directly solving the primal prob-lems sometimes is very difficult; however, in this paper, their prime-dual Gao-Strang
complementary problems enable us to elegantly and easily not only get the global optimal solutions of primal problems but also of canonical dual problems. The canonical duality theory (CDT) was originally developed for nonconvex/nonsmooth mechanics [7]. It is now realized that this potentially powerful theory can be used for solving a large class of nonconvex/nonsmooth GO problems [3, 9, 10, 12, 13, 14, 15] with applications to data mining clustering, sensor network problems, and molecular distance geometry problem [20], etc.. In this paper, a CDT-unified GO algorithm is presented. According to the CDT, the proof of the convergent theorem for the algorithm is very easy and clear. High-dimensional nonconvex/nonsmooth GO examples (such as the minimization problem of Rosenbrock function) are tested by the new algorithm and numerical results pleasantly show that the new algorithm has a great promise for solving some GO problems.

## 2 The Algorithm and its Convergence

The algorithm is established from CDT [7]. In this paper we will solve the following nonconvex GO problem [18, 16]:

$$
\begin{equation*}
(\mathcal{P}) \quad \min _{x}\left\{P(x)=V(\Lambda(x))-F(x) \mid x \in \mathcal{X}_{a}\right\}, \tag{1}
\end{equation*}
$$

where $\mathcal{X}_{a} \subset \mathbb{R}^{n}$ is a feasible space, $V(\Lambda(x))$ is not necessarily convex with respect to $x$ and it is a so-called canonical function satisfying

$$
\begin{align*}
& V^{*}(\varsigma)=\operatorname{sta}\{\langle\xi ; \varsigma\rangle-V(\xi) \mid \xi \in \mathcal{V}\}: \mathcal{V}^{*} \rightarrow \mathbb{R},  \tag{2}\\
& \varsigma=\nabla V(\xi) \Leftrightarrow \xi \in \nabla V^{*}(\varsigma) \Leftrightarrow V(\xi)+V^{*}(\varsigma)=\langle\xi ; \varsigma\rangle, \tag{3}
\end{align*}
$$

$\Lambda$ is a geometrically admissible (objective) mapping from the feasible space $\mathcal{X}_{a}$ into a canonical measure space $\mathcal{V}, F(x)=\langle x, f\rangle-\frac{1}{2}\langle x, A x\rangle,\langle *, *\rangle$ denotes a bilinear form in $\mathbb{R}^{n} \times \mathbb{R}^{n}, f \in \mathbb{R}^{n}, A=A^{T} \in \mathbb{R}^{n \times n}$ are given, and $\langle * ; *\rangle$ represents a bilinear form which puts $\mathcal{V}$ and $\mathcal{V}^{*}$ in duality. By (22)-(3), the nonconvex function $P(x)$ in (11) can be rewritten as

$$
\begin{equation*}
\Xi(x, \varsigma)=\langle\Lambda(x) ; \varsigma\rangle-V^{*}(\varsigma)-F(x) . \tag{4}
\end{equation*}
$$

In real applications, the objective operator $\Lambda$ is usually a quadratic measure over a given a field [7, 18] and in finite space this quadratic measure can be written as [8, 18, 16]:

$$
\begin{equation*}
\Lambda(x)=\left\{\frac{1}{2} x^{T} B_{k} x+b_{k}^{T} x+c_{k}=\xi_{k}, k=1,2, \ldots, m\right\}=\xi: \mathbb{R}^{n} \rightarrow \mathbb{R}^{m} \tag{5}
\end{equation*}
$$

where $B_{k} \in \mathbb{R}^{n \times n}$ is a symmetrical matrix, $b_{k} \in \mathbb{R}^{n}, c_{k} \in \mathbb{R}^{1}$ for each $k=1,2, \ldots, m$. Denote

$$
\begin{align*}
& G(\varsigma)=A+\sum_{k=1}^{m} \varsigma_{k} B_{k}: \mathbb{R}^{m} \rightarrow \mathbb{R}^{n \times n},  \tag{6}\\
& d(\varsigma)=f-\sum_{k=1}^{m} \varsigma_{k} b_{k}: \mathbb{R}^{m} \rightarrow \mathbb{R}^{n}, \tag{7}
\end{align*}
$$

and define

$$
\begin{align*}
& S_{a}=\left\{\varsigma \in \mathbb{R}^{m} \mid d(\varsigma) \in \operatorname{Col}(G(\varsigma))\right\},  \tag{8}\\
& S_{a}^{+}=\left\{\varsigma \in S_{a} \mid G(\varsigma) \succeq 0\right\} \tag{9}
\end{align*}
$$

where $\operatorname{Col}(G(\varsigma))$ is the space spanned by the column of $G(\varsigma)$, and $G(\varsigma) \succeq 0$ stands for $G(\varsigma)$ is a semi-positive definite matrix (if $\operatorname{det} G(\varsigma)=0, G(\varsigma)^{-1}$ should be understood as the generalized inverse [8, 18, 16]).

The following algorithm is designed:
Algorithm 1 - A canonical dual theory algorithm.
Step 1. Call the subroutine Algorithm 2, or simply the Matlab's fsolve program if dimension of problems are less than a few thousands, to solve differential equations $\Xi(x, \varsigma)^{\prime}=0$.

Step 2. Output the roots $\bar{\varsigma}, \bar{x}$.
Step 3. Check whether $\bar{\varsigma} \in S_{a}$ are in $S_{a}^{+}$; if so, then
Step 4. Pick up the $\bar{x}$ s that are satisfying $G(\bar{\zeta}) \bar{x}=d(\bar{\zeta})$.

Theorem 1 (Convergence of the Algorithm) $\bar{x}$ generated by Algorithm 1 is the global optimal solution of $(\mathcal{P})$.

Proof. By Theorem 3 of [18] and its reference [6] we know that $\bar{x}$ generated by Algorithm 1 is a global optimal solution of $(\mathcal{P})$.

The powerful of this simple algorithm can be easily demonstrated by easily and effectively solving the following benchmark test problems of GO, even calculated by hands on paper.

Example 1. (Two-dimensional Rosenbrock function) $P(x)=\left(x_{1}-1\right)^{2}+100\left(x_{2}-\right.$ $\left.x_{1}^{2}\right)^{2}$.
Solution. $\Xi(x, \varsigma)=\left(x_{1}-1\right)^{2}+\left(x_{1}^{2}-x_{2}\right) \varsigma-\frac{1}{400} \varsigma^{2}$ and $S_{a}^{+}=\left\{\varsigma \in \mathbb{R}^{1} \mid \varsigma>-1\right\}$ are got. Let $\Xi(x, \varsigma)^{\prime}=0$ a critical point $\left(\bar{x}_{1}, \bar{x}_{2}, \bar{\varsigma}\right)=(1,1,0)$ is got. We easily know $\bar{\varsigma}=0 \in S_{a}^{+}$ and satisfying $G(\bar{\varsigma}) \bar{x}=\left(\begin{array}{ll}2 & 0 \\ 0 & 2\end{array}\right)\binom{1}{1}=\binom{2}{2}=f=d(\bar{\varsigma})$. Thus, $\left(\bar{x}_{1}, \bar{x}_{2}\right)=(1,1)$ is a global minimum of Rosenbrock function.

Example 2. (Two-dimensional De Jong function) $P(x)=100\left(x_{2}-x_{1}^{2}\right)^{2}+\left(x_{1}-\right.$ $1)^{2}+\left(x_{2}-1\right)^{2}$.
Solution. $\Xi(x, \varsigma)=\left(x_{1}-1\right)^{2}+\left(x_{2}-1\right)^{2}+\left(x_{1}^{2}-x_{2}\right) \varsigma-\frac{1}{400} \varsigma^{2}$ and $S_{a}^{+}=\left\{\varsigma \in \mathbb{R}^{1} \mid \varsigma>-1\right\}$ are got. Let $\Xi(x, \varsigma)^{\prime}=0$ a critical point $\left(\bar{x}_{1}, \bar{x}_{2}, \bar{\varsigma}\right)=(1,1,0)$ is got. We easily know $\bar{\varsigma}=0 \in S_{a}^{+}$and satisfying $G(\bar{\varsigma}) \bar{x}=f=d(\bar{\varsigma})$. Thus, $\left(\bar{x}_{1}, \bar{x}_{2}\right)=(1,1)$ is a global minimal
solution.
Example 3. (Colville function) $P(x)=100\left(x_{2}-x_{1}^{2}\right)^{2}+90\left(x_{3}^{2}-x_{4}\right)^{2}+\left(x_{1}-1\right)^{2}+$ $\left(x_{3}-1\right)^{2}+10.1\left(\left(x_{2}-1\right)^{2}+\left(x_{4}-1\right)^{2}\right)+19.8\left(x_{2}-1\right)\left(x_{4}-1\right)$.
Solution. Let $\Xi(x, \varsigma)^{\prime}=0$ and a unique critical point $\left(\bar{x}_{1}, \bar{x}_{2}, \bar{x}_{3}, \bar{x}_{4}, \bar{\varsigma}_{1}, \bar{\varsigma}_{2}\right)=(1,1,1,1,0,0)$ is got such that $\left(\bar{\varsigma}_{1}, \bar{\varsigma}_{2}\right) \in S_{a}^{+}$and $G(\bar{\varsigma}) \bar{x}=f=d(\bar{\varsigma})$. Thus, $\left(\bar{x}_{1}, \bar{x}_{2}, \bar{x}_{3}, \bar{x}_{4}\right)=(1,1,1,1)$ is a global minimal solution.

Example 4. (Two-dimensional Zakharov function) $P(x)=x_{1}^{2}+x_{2}^{2}+\left(0.5 x_{1}+\right.$ $\left.x_{2}\right)^{2}+\left(0.5 x_{1}+x_{2}\right)^{4}$.
Solution. Let $\Xi(x, \varsigma)^{\prime}=0$ and a unique critical point $\left(\bar{x}_{1}, \bar{x}_{2}, \bar{\varsigma}\right)=(0,0,0)$ is got such that $\bar{\varsigma} \in S_{a}^{+}$, satisfying the formula $G(\bar{\varsigma}) \bar{x}=\left(\begin{array}{cc}5 / 2 & 1 \\ 1 & 4\end{array}\right)\binom{0}{0}=\binom{0}{0}=f=d(\bar{\varsigma})$. Thus, $(0,0)$ is the global optimal solution.

Example 5. (Four-dimensional Powell function) $P(x)=\left(x_{1}-10 x_{2}\right)^{2}+5\left(x_{3}-\right.$ $\left.x_{4}\right)^{2}+\left(x_{2}-x_{3}\right)^{4}+10\left(x_{1}-x_{4}\right)^{4}$.
Solution. Let $\Xi(x, \varsigma)^{\prime}=0$ and a unique critical point $\left(\bar{x}_{1}, \bar{x}_{2}, \bar{x}_{3}, \bar{x}_{4}, \bar{\varsigma}_{1}, \bar{\varsigma}_{2}\right)=(0,0,0,0,0,0)$ is got such that $\left(\bar{\varsigma}_{1}, \bar{\varsigma}_{2}\right) \in S_{a}^{+}$and $G(\bar{\varsigma}) \bar{x}=f=d(\bar{\varsigma})$. Thus, $\left(\bar{x}_{1}, \bar{x}_{2}, \bar{x}_{3}, \bar{x}_{4}\right)=(0,0,0,0)$ is a global minimal solution.

Example 6. (Double Well function) $P(x)=\frac{1}{2}\left(\frac{1}{2} x^{2}-2\right)^{2}-\frac{1}{2} x$.
Solution. $\Xi(x, \varsigma)=\left(\frac{1}{2} x^{2}-2\right) \varsigma-\frac{1}{2} \varsigma^{2}-\frac{1}{2} x$ and $S_{a}^{+}=\left\{\varsigma \in \mathbb{R}^{1} \mid \varsigma>0\right\}$ are got. Let $\Xi(x, \varsigma)^{\prime}=0$ and three critical points of $\Xi(x, \varsigma):\left(\bar{x}^{1}, \bar{\varsigma}^{1}\right)=(2.11491,0.236417),\left(\bar{x}^{2}, \bar{\varsigma}^{2}\right)=$ $(-1.86081,-0.268701),\left(\bar{x}^{3}, \bar{\varsigma}^{3}\right)=(-0.254102,-1.96772)$ are got, and a unique point $\bar{\varsigma}^{1}$ in $S_{a}^{+}$is got. Thus, $\bar{x}^{1}=2.11491$ is the global minimal solution of the Double Well function. This can also be illuminated in Figure 1 .


Figure 1: The prime double-well function.
The canonical dual problem of Double Well function is:

$$
\begin{equation*}
\max _{\varsigma>0}-\frac{1}{8 \varsigma}-\frac{1}{2} \varsigma^{2}-2 \varsigma, \tag{10}
\end{equation*}
$$

which is equal to

$$
\begin{equation*}
\min _{\varsigma>0} \frac{1}{8 \varsigma}+\frac{1}{2} \varsigma^{2}+2 \varsigma \tag{11}
\end{equation*}
$$

that can be transferred into (but not equal to (because of the constant of $\varsigma_{0}$ )) (this idea is communicated from Wu C.Z.)

$$
\min _{t, \varsigma} t \quad \text { subject to }\left(\begin{array}{cc}
\varsigma & 0.5 / \sqrt{2}  \tag{12}\\
0.5 / \sqrt{2} & t-\frac{1}{2} \varsigma_{0}^{T} \varsigma-2 \varsigma
\end{array}\right) \succeq 0
$$

and solved by Semi-Definite Programming (SDP) package SeDuMi 1_3 [24] (this idea is communicated from Wu C.Z. In this whole paper, this is the single idea of other researchers). However, after 6 circles in all 33 iterations SeDuMi 1_3 died at the solution 0.2585 , which is still very far from the real global optimal solution 0.236417 of the dual problem (10). This means the popular SDP package SeDuMi 1_3 has a poor performance for Double Well function minimization problem, compared with our Algorithm 1. But other strategies to solve the following dual problem of (1) are sought:

$$
\begin{equation*}
\min _{G(\varsigma) \succeq 0} P^{d}(\varsigma)=\frac{1}{2} d(\varsigma)^{T} G(\varsigma)^{-1} d(\varsigma)+\sum_{k=1}^{m}\left(\frac{1}{2} \alpha_{k}^{-1} \varsigma_{k}^{2}-c_{k} \varsigma_{k}\right) ; \tag{13}
\end{equation*}
$$

for example, one strategy is to replace (13) by the following quadratic semidefinite programming problem (QSDP)

$$
\min _{t, \varsigma} \quad t+\frac{1}{2} \alpha^{-1} \varsigma^{T} \varsigma++\varsigma^{T} c \quad \text { subject } \quad \text { to }\left(\begin{array}{cc}
G(\varsigma) & \frac{1}{\sqrt{2}} d(\varsigma)  \tag{14}\\
\frac{1}{\sqrt{2}} d(\varsigma)^{T} & t
\end{array}\right) \succeq 0,
$$

where $\alpha, c, \varsigma$ are vectors, and some known QSDP packages can solve (14) for lowdimensional problems and for high-dimensional problems we can design QSDP algorithms by ourselves. This is one direction of algorithm design for CDT. Another direction is to efficiently and effectively get the roots of $\Xi(x, \varsigma)^{\prime}=0$ for Algorithm 1, i.e. to solve the $m+n$ quadratic equations (15) given as follows.

For high-dimensional nonconvex GO problems, the finite element method (FEM)based [25] subroutine of Algorithm 1 is well designed as follows. Step 1 of Algorithm 1 is to find the roots of $\Xi(x, \varsigma)^{\prime}=0$, i.e. the following $m+n$ quadratic equations:

$$
\begin{cases}\frac{1}{2} x^{T} B_{k} x+b_{k}^{T} x+c_{k}=\alpha_{k}^{-1} \varsigma_{k}, k=1,2, \ldots, m & \left(b y \quad \Xi(x, \varsigma)_{\varsigma}^{\prime}=0\right),  \tag{15}\\ G(\varsigma) x=d(\varsigma), \quad \text { i.e. } \quad\left(A+\sum_{k=1}^{m} \varsigma_{k} B_{k}\right) x=f-\sum_{k=1}^{m} \varsigma_{k} b_{k} & \left(\text { by } \quad \Xi(x, \varsigma)_{x}^{\prime}=0\right),\end{cases}
$$

where $\alpha_{k}, k=1,2, \ldots, m$ are the coefficients in $V(\Lambda(x))=\sum_{k=1}^{m} \frac{1}{2} \alpha_{k}\left(\frac{1}{2} x^{T} B_{k} x+b_{k}^{T} x+\right.$ $\left.c_{k}\right)^{2}$.

Algorithm 2-A subroutine finding roots of (15). E.g. the finite element discretized $\Xi(x, \varsigma)^{\prime}$ method if the dimension of the problem is large than a few thousands.

Firstly the minimization problem for Rosenbrock function for (15) is tested. In global optimization, the nonconvex minimization problem of Rosenbrock function is a benchmark test problem that is extensively used to test the performance of optimization algorithms and approaches. The global minimum is inside a long, deep, narrow, parabolic/banana shaped flat valley. The shallow global minimum is inside a deeply curved valley. To find the valley and to converge to the global minimum is difficult. Detailed introduction to the difficulty of the problem and the excellence as the testing problem of GO can be referred to [19].

Minimizing the Rosenbrock function:

$$
\begin{align*}
\min P(x) & =\sum_{k=1}^{n-1}\left[100\left(x_{k}^{2}-x_{k+1}\right)^{2}+\left(x_{k}-1\right)^{2}\right]: x \in \mathbb{R}^{n}  \tag{16}\\
& =n-1+\sum_{k=1}^{n-1} \frac{1}{2} \alpha_{k}\left(x_{k}^{2}-x_{k+1}\right)^{2}+\sum_{k=1}^{n-1} x_{k}^{2}-\sum_{k=1}^{n-1} 2 x_{k}  \tag{17}\\
& =n-1+\sum_{k=1}^{n-1} \frac{1}{2} \alpha_{k}\left(\frac{1}{2} x^{T} B_{k} x+b_{k}^{T} x+c_{k}\right)^{2}+\frac{1}{2} x^{T} A x-x^{T} f, \tag{18}
\end{align*}
$$


$S_{a}^{+}=\left\{\varsigma_{k}>-1, k=1,2, \ldots, n-1\right\}$. For Rosenbrock function, (15) is written as:

$$
\left\{\begin{array}{l}
x_{k}^{2}-x_{k+1}=0.005 \varsigma_{k}, k=1,2, \ldots, n-1  \tag{19}\\
2\left(1+\varsigma_{k+1}\right) x_{k+1}=2+\varsigma_{k}, k=1,2, \ldots, n-3 \\
\varsigma_{n-1}=0 \\
\left(1+\varsigma_{1}\right) x_{1}=1, \\
2 x_{n-1}=2+\varsigma_{n-2}
\end{array}\right.
$$

(19) is solved by Matlab's $f$ solve on the $\operatorname{Intel}(\mathrm{R})$ Celeron(R) CPU $900 @ 2.20 \mathrm{GHz}$ Windows Vista ${ }^{\mathrm{TM}}$ Home Basic personal notebook computer. The initial solution for ( $x ; \varsigma$ ) is set as $(3, \ldots, 3 ; 2, \ldots, 2)$ and the numerical computational results are shown in Table 1. We may see in Table 1 that the global minimal solution $(1, \ldots, 1)$ for (1) can be easily and accurately got by Algorithm 1, at the same time the global maximal solution $(0, \ldots, 0)$ for the (13) over $S_{a}^{+}$can be easily and accurately got by Algorithm 1 directly too. The smart Matlab's fsolve does not output any other solution in $S / S_{a}^{+}$.

Table 1: Results of numerical computations

| Dimension $n$ | Iterations | Calls of Equations (19) | prime and dual opt sln |  | CPU time |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 3 | 6 | 42 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 0.0468 |
| 4 | 6 | 56 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 0.0468 |
| 5 | 6 | 70 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 0.0624 |
| 6 | 7 | 96 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 0.0780 |
| 7 | 7 | 112 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 0.1248 |
| 8 | 7 | 128 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 0.1248 |
| 9 | 7 | 144 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 0.0936 |
| 10 | 7 | 160 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 0.1404 |
| 20 | 7 | 320 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 0.2808 |
| 30 | 8 | 540 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 0.4368 |
| 40 | 8 | 720 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 0.6084 |
| 50 | 8 | 900 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 0.8736 |
| 60 | 8 | 1,080 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 1.0920 |
| 70 | 8 | 1,260 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 1.2948 |
| 80 | 8 | 1,440 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 1.6536 |
| 90 | 8 | 1,620 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 1.9344 |
| 100 | 8 | 1,800 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 2.2464 |
| 200 | 9 | 4,000 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 7.7844 |
| 300 | 9 | 6,000 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 15.8497 |
| 400 | 9 | 8,000 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 27.4250 |
| 500 | 9 | 10,000 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 42.4947 |
| 600 | 9 | 12,000 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 61.2616 |
| 700 | 9 | 14,000 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 88.2342 |
| 800 | 9 | 16,000 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 118.3736 |
| 900 | 10 | 19,800 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 170.8367 |
| 1,000 | 10 | 22,000 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 221.8334 |
| 2,000 | 10 | 44,000 | $(1, \ldots, 1)$ | $(0, \ldots, 0)$ | 1491.9 |

Chemical database clustering problem. Algorithm 1 is applied to solve the real chemical database clustering problem (34)-(35) of [27], i.e.

$$
\begin{align*}
\min P\left(Y_{1}, Y_{2}, \ldots, Y_{n}\right) & =\sum_{i=1}^{n-1} \sum_{j=i+1}^{n} \frac{1}{2} \alpha_{i j}\left(\left\|Y_{i}-Y_{j}\right\|^{2}-d_{i j}^{2}\right)^{2}  \tag{20}\\
& =\sum_{i, j=1, i<j}^{n} \frac{1}{2} \alpha_{i j}\left(\left(y_{i 1}-y_{j 1}\right)^{2}+\left(y_{i 2}-y_{j 2}\right)^{2}-d_{i j}^{2}\right)^{2}, \tag{21}
\end{align*}
$$

where $\alpha_{i j}=\frac{1}{2 d_{i j}^{4}}$ if $d_{i j}^{4} \geq 10^{-12}, \alpha_{i j}=\frac{1}{2}$ if $d_{i j}^{4}<10^{-12}$, for all $i, j=1,2, \ldots, n\left(d_{i j}\right.$ is calculated from the original $n$ by 9 dataset [27]) and $Y_{i} \in \mathbb{R}^{2}, i=1,2, \ldots, n$. The Gao-Strang generalized complementary function for (21) is:

$$
\begin{align*}
\Xi(Y, \varsigma) & =\sum_{i, j=1, i<j}^{n}\left(\left(\left\|Y_{i}-Y_{j}\right\|^{2}-d_{i j}^{2}\right) \varsigma_{i j}-\frac{1}{2 \alpha_{i j}} \varsigma_{i j}^{2}\right)  \tag{22}\\
& =\sum_{i, j=1, i<j}^{n}\left(\left(\left(y_{i 1}-y_{j 1}\right)^{2}+\left(y_{i 2}-y_{j 2}\right)^{2}-d_{i j}^{2}\right) \varsigma_{i j}-\frac{1}{2 \alpha_{i j}} \varsigma_{i j}^{2}\right) . \tag{23}
\end{align*}
$$

Choose the SVD-reduced initial solution $Y_{i}^{0}, i=1,2, \ldots, n$ [27, 29] and calculate $\varsigma_{i j}^{0}=$ $\alpha_{i j}\left(\left\|Y_{i}^{0}-Y_{j}^{0}\right\|^{2}-d_{i j}^{2}\right)$ as the initial solution for $\varsigma_{i j}^{0}$ in solving the following quadratic nonlinear equations of $\Xi(Y, \varsigma)^{\prime}=0$ :

$$
\left\{\begin{array}{l}
2\left(y_{i 1}-y_{j 1}\right) \varsigma_{i j}=0, i=1,2, \ldots, n-1, j=i+1, \ldots, n  \tag{24}\\
-2\left(y_{i 1}-y_{j 1}\right) \varsigma_{i j}=0, i=1,2, \ldots, n-1, j=i+1, \ldots, n \\
2\left(y_{i 2}-y_{j 2}\right) \varsigma_{i j}=0, i=1,2, \ldots, n-1, j=i+1, \ldots, n \\
-2\left(y_{i 2}-y_{j 2}\right) \varsigma_{i j}=0, i=1,2, \ldots, n-1, j=i+1, \ldots, n \\
\left(y_{i 1}-y_{j 1}\right)^{2}+\left(y_{i 2}-y_{j 2}\right)^{2}-d_{i j}^{2}-\frac{1}{\alpha_{i j}} \varsigma_{i j}=0, i=1,2, \ldots, n-1, j=i+1, \ldots, n
\end{array}\right.
$$

The global optimal prime or dual solution is got and the comparison of Algorithm 1 with the SD (steepest descent) method, CG (conjugate gradient) method, BFGS (approximated Newton) method, NR (Newton-Raphson) method, and TN (truncated Newton) method T-IHN (truncated incomplete Hessian Newton) method, etc can be made. The equations (24) are illuminated to make readers to easily understand the equations. When $n=2$, i.e. $i=1, j=2$ there are the following 5 quadratic equations with 5 variables (with initial value $\left(y_{11}^{0}, y_{12}^{0}, y_{21}^{0}, y_{22}^{0} ; \varsigma_{12}^{0}\right)=(331.5590,-188.4908,364.6889,-158.1010 ; 0)$ ):

$$
\left\{\begin{array}{l}
2\left(y_{11}-y_{21}\right) \varsigma_{12}=0  \tag{25}\\
-2\left(y_{11}-y_{21}\right) \varsigma_{12}=0 \\
2\left(y_{12}-y_{22}\right) \varsigma_{12}=0 \\
-2\left(y_{12}-y_{22}\right) \varsigma_{12}=0 \\
\left(y_{11}-y_{21}\right)^{2}+\left(y_{12}-y_{22}\right)^{2}-15701874.4205486 \varsigma_{12}-2801.95239257813=0
\end{array}\right.
$$

In (24), different database of [29] has different $n$. Numerical computational results of solving (24) will be updated.

For CDT, there are two research directions for its algorithm design. One is to design the CDT algorithm to solve (13); for example, one strategy is to design the quadratic semidefinite programming (QSDP) algorithm to solve (14). Another research direction is to design the CDT algorithm to solve the special $\mathrm{m}+\mathrm{n}$ quadratic (nonlinear) equations, with $\mathrm{m}+\mathrm{n}$ variables, (15); for example, Newton-type iteration algorithms, gradient-type iteration algorithms, trust-region-type iteration algorithms, nonlinear finite-element-type algorithms, Hamiltonian system symplectic-type algorithms, etc are good strategy to solve (15). Researchers may design a powerful CDT algorithm along these two directions to solve (13) and (15) respectively.

## 3 Advantages of Solving the Dual Problem Than Solving the Prime Problem

In this section, we still use the benchmark GO test problem, minimizing the Rosenbrock function, to illuminate some advantages of solving the canonical dual problem (13) compared with directly solving the prime problem (1). By the CDT [7, 16, the canonical dual problem of (16) is:

$$
\begin{equation*}
\max _{\varsigma>-1} P^{d}(\varsigma)=n-1-\sum_{k=1}^{n-1}\left[\frac{\left(\varsigma_{k-1}+2\right)^{2}}{4\left(\varsigma_{k}+1\right)}+\frac{1}{400} \varsigma_{k}^{2}\right], \tag{26}
\end{equation*}
$$

where $\varsigma_{0}=0$. (26) is solved by the Discrete Gradient (DG) method [1], a local search optimization solver for nonconvex and/or nonsmooth optimization problems, and the numerical computational results are listed in Tables 2/3 (where seed1 is the initial solution $\left(x^{0} ; \varsigma^{0}\right)=(3,3, \ldots, 3 ;-2 / 3,-2 / 3, \ldots,-2 / 3,0)$ and seed2 is the initial solution $\left.\left(x^{0} ; \varsigma^{0}\right)=(100,100, \ldots, 100 ; 100,100, \ldots, 100,0)\right)$.

In Tables 233, we may see that the dual problem (26) can be elegantly, easily, quickly and accurately solved to get its objective function value 0.00000000 , compared with the prime problem (16). We may also see that (26) is convenient for MPI (Message Passing Interface) parallel computation. The successfully tested MPI code is followed:

```
broadcast \(n-1\)
    call MPI_BCAST ( \(n-1,1\), MPI_INTEGER, 0, MPI_COMM_WORLD , ierr)
check for quit signal
    if ( \(n-1\).le. 0 ) goto 30
calculate every partials
    sum \(=0.0 \mathrm{~d} 0\)
    do \(20 \mathrm{i}=\) myid \(+1, n-1\), numprocs
        if \((i-1\).eq. 0\()\) then \(\varsigma(0)=0\)
        sum \(=\operatorname{sum}+(\varsigma(i-1)+2.0) /(4 *(\varsigma(i)+1.0))+(1.0 / 400.0) * \varsigma(i) * * 2\)
    20 continue
    \(\mathrm{f}=\mathrm{sum}\)
```

collect all the partial sums
call MPI_REDUCE (f,objf,1,MPI_DOUBLE_PRECISION, MPI_SUM, 0 ,
\& MPI_COMM_WORLD, ierr )

30 node 0 (i.e. myid $=0$ ) prints the sums $=$ objf

Table 2: Results of numerical experiments for seed1

| Dimension $n$ | Iterations |  | Function calls |  | Objective function value |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Prime | Dual | Prime | Dual | Prime | Dual |
| 2 | 120 | 24 | 2843 | 28 | 0.00001073 | 0.00000000 |
| 3 | 422 | 26 | 8996 | 137 | 0.00401438 | 0.00000000 |
| 4 | 3737 | 35 | 48352 | 202 | 0.00615273 | 0.00000000 |
| $5^{*}$ | 335 | 34 | 10179 | 399 | 3.96077434 | 0.00000000 |
| $6^{*}$ | 2375 | 44 | 43770 | 868 | 4.00635895 | 0.00000000 |
| $7^{*}$ | 1223 | 53 | 28009 | 1625 | 4.09419146 | 0.00000000 |
| 8 | 2160 | 55 | 46792 | 2100 | 0.01246714 | 0.00000000 |
| 9 | 2692 | 51 | 61017 | 2526 | 0.01397307 | 0.00000000 |
| 10 | 4444 | 63 | 91470 | 3979 | 0.01055630 | 0.00000000 |
| 20 | 3042 | 55 | 140924 | 10084 | 0.00940077 | 0.00000000 |
| 30 | 2321 | 58 | 133980 | 20515 | 0.01075478 | 0.00000000 |
| 40 | 1659 | 60 | 173795 | 26818 | 0.01227866 | 0.00000000 |
| 50 | 2032 | 57 | 219233 | 36459 | 0.01264147 | 0.00000000 |
| 60 | 1966 | 61 | 260701 | 50495 | 0.01048188 | 0.00000000 |
| 70 | 1876 | 56 | 272919 | 52545 | 0.01531147 | 0.00000000 |
| 80 | 1405 | 61 | 195156 | 59684 | 0.01594730 | 0.00000000 |
| 90 | 2142 | 61 | 371963 | 71320 | 0.01055831 | 0.00000000 |
| 100 | 2676 | 60 | 510722 | 70208 | 0.01125514 | 0.00000000 |
| 200 | 1395 | 61 | 653604 | 188589 | 0.01115318 | 0.00000000 |
| 300 | 1368 | 60 | 882760 | 235163 | 0.01574873 | 0.00000000 |
| 400 | 2085 | 66 | 1869675 | 301805 | 0.00928066 | 0.00000000 |
| 500 | 1155 | 59 | 1394240 | 358938 | 0.01168440 | 0.00000000 |
| 600 | 1226 | 63 | 1808285 | 451817 | 0.00918730 | 0.00000000 |
| 700 | 1557 | 60 | 2134359 | 559378 | 0.01257100 | 0.00000000 |
| 800 | 1398 | 61 | 2098062 | 522726 | 0.01442714 | 0.00000000 |
| 900 | 716 | 65 | 1904187 | 763449 | 0.01074534 | 0.00000000 |
| 1000 | 1825 | 61 | 3598608 | 681509 | 0.00897202 | 0.00000000 |
| 2000 | 257 | 62 | 2087277 | 1455472 | 0.00937219 | 0.00000000 |
| 3000 | 3221 | 60 | 20642543 | 2714296 | 0.01250373 | 0.00000000 |
| $4000^{*}$ | 679 | 60 | 7581502 | 3659292 | 4.11193171 | 0.00000000 |

This shows the advantages of solving the dual problem (13) than solving the prime problem (11). We may use maximizing the canonical dual function of Colville function (see Example 3) to illuminate this point again.

It is known that directly solving the minimizing of Colville function is difficult. But, it is very easy to solving the following canonical dual problem:

$$
\begin{aligned}
& \max _{\varsigma>-1} P^{d}(\varsigma)=42-\frac{1}{400} \varsigma_{1}^{2}-\frac{1}{360} \varsigma_{2}^{2} \\
& -\frac{1}{2}\left(\begin{array}{c}
2 \\
40+\varsigma_{1} \\
2 \\
40+\varsigma_{2}
\end{array}\right)^{T}\left(\begin{array}{cccc}
2+2 \varsigma_{1} & 0 & 0 & 0 \\
0 & 20.2 & 0 & 19.8 \\
0 & 0 & 2+2 \varsigma_{2} & 0 \\
0 & 19.8 & 0 & 20.2
\end{array}\right)+\left(\begin{array}{c}
2 \\
40+\varsigma_{1} \\
2 \\
40+\varsigma_{2}
\end{array}\right)
\end{aligned}
$$

Table 3: Results of numerical experiments for seed2

| Dimension $n$ | Iterations |  | Function calls |  | Objective function value |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Prime | Dual | Prime | Dual | Prime | Dual |
| 2 | 10013 | 24 | 227521 | 28 | 47.23824896 | 0.00000000 |
| 3 | 144 | 32 | 4869 | 235 | 96.49814330 | 0.00000000 |
| 4 | 144 | 81 | 5279 | 938 | 82.46230602 | 0.00000000 |
| 5 | 148 | 137 | 5682 | 1768 | 94.19254867 | 0.00000000 |
| 6 | 154 | $166^{*}$ | 6238 | 2590 | 88.84382963 | 0.00000000 |
| 7 | 159 | $179^{*}$ | 7097 | 3288 | 237.63078399 | 0.00000000 |
| 8 | 165 | $202^{*}$ | 7502 | 4300 | 238.41126013 | 0.00000000 |
| 9 | 153 | $206^{*}$ | 7137 | 5083 | 84.54205412 | 0.00000000 |
| 10 | 162 | $216^{*}$ | 7491 | 5920 | 83.23094398 | 0.00000000 |
| 20 | 225 | $285^{*}$ | 19111 | 17458 | 83.94779152 | 0.00000000 |
| 30 | 216 | $301^{*}$ | 20939 | $28543^{*}$ | 156.95838274 | 0.00000000 |
| 40 | 163 | $291^{*}$ | 19775 | $40444^{*}$ | 83.30960344 | 0.00000000 |
| 50 | 158 | $298^{*}$ | 33269 | $51888^{*}$ | 85.93091895 | 0.00000000 |
| 60 | 158 | $312^{*}$ | 34094 | $61767^{*}$ | 89.07412094 | 0.00000000 |
| 70 | 162 | $284^{*}$ | 35436 | $69865^{*}$ | 92.45725362 | 0.00000000 |
| 80 | 209 | $297^{*}$ | 35607 | $89127^{*}$ | 157.69955825 | 0.00000000 |
| 90 | 227 | $294^{*}$ | 60398 | $98748^{*}$ | 82.44035053 | 0.00000000 |
| 100 | 202 | $290^{*}$ | 57792 | $102796^{*}$ | 81.94595276 | 0.00000000 |
| 200 | 1826 | 262 | 436413 | 189293 | 83.77165551 | 0.00000000 |
| 300 | 195 | $259^{*}$ | 169238 | $261320^{*}$ | 152.95671738 | 0.00000000 |
| 400 | 195 | $278^{*}$ | 212104 | $375816^{*}$ | 82.49253919 | 0.00000000 |
| 500 | 190 | $297^{*}$ | 331637 | $52265^{*}$ | 82.40170647 | 0.00000000 |
| 600 | 292 | $303^{*}$ | 431092 | $559068^{*}$ | 150.15456693 | 0.00000000 |
| 700 | 189 | $275^{*}$ | 383735 | $758633^{*}$ | 89.14575473 | 0.00000000 |
| 800 | 198 | $270^{*}$ | 429674 | $701053^{*}$ | 84.50538257 | 0.00000000 |
| 900 | 198 | $280^{*}$ | 416150 | $867398^{*}$ | 85.32757049 | 0.00000000 |
| 1000 | 193 | $283^{*}$ | 445326 | $930761^{*}$ | 89.48369379 | 0.00000000 |
| 2000 | 232 | $310^{*}$ | 1123240 | $2030104^{*}$ | 84.26810981 | 0.00000000 |

which is a simple problem of maximizing a concave function over a convex set: by watching Figure 2, $\bar{\varsigma}=(0,0)$ is easily known as the global optimal solution.

Thus, the following optimization algorithm designed to solve the canonical dual problem (14) is very necessary.

Algorithm 3-An optimization algorithm to solve the Quadratic Semidefinite Programming (14).

In Example 6, (11) is equal to

$$
\begin{aligned}
& \min _{t, \varsigma \in \mathbb{R}} \frac{1}{2}\binom{t}{\varsigma}^{T}\left(\begin{array}{cc}
0 & 0 \\
0 & 1
\end{array}\right)\binom{t}{\varsigma}+\binom{1}{2}^{T}\binom{t}{\varsigma} \\
& \text { s.t. } \quad\left(\begin{array}{cc}
\varsigma & 0.5 / \sqrt{2} \\
0.5 / \sqrt{2} & t
\end{array}\right) \succeq 0
\end{aligned}
$$



Figure 2: Canonical dual Colville function on $S_{a}^{+}=\{\varsigma>-1\}$
which can solved by Algorithm 3.

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